

A MODEL OF "NONADDITIVE" ROUTING PROBLEM WHERE THE COSTS DEPEND ON THE SET OF PENDING TASKS

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We consider a generalization of the bottleneck (minimax) routing problem. The problem is to successively visit a number of megalopolises, complicated by precedence of constraints imposed on the order of megalopolises visited and the fact that the cost functions (of movement between megalopolises and of interior tasks) may explicitly depend on the list of tasks that are not completed at the present time. The process of movement is considered to be a sequence of steps, which include the exterior movement to the respective megalopolis and the following completion of (essentially interior) jobs connected with the megalopolis. The quality of the whole process is represented by the maximum cost of steps it consists of; the problem is to minimize the mentioned criterion (which yields a minimax problem, usually referred to as a "bottleneck problem"). Optimal solutions, in the form of a route-track pair (a track, or trajectory, conforms to a specific instance of a tour over the megalopolises, which are numbered in accordance with the route; the latter is defined by the transposition of indices), are constructed through a "nonstandard" variant of the dynamic programming method, which allows to avoid the process of constructing of all the values of the Bellman function whenever precedence constraints are present.

Keywords: dynamic programming; route; precedence constraints; sequential ordering problem.

Introduction

The paper investigates the problem of sequentially walking through a set of megalopolises, where the cost is aggregated in a nonadditive way and the walk is restricted by precedence constraints. The cost functions, which measure the quality of the steps of the process, depend on the list of "pending" tasks. The mentioned dependence may arise in real-world problems connected with movement through a radiation field, for example, while dismantling ionizing radiation sources. Other applications are possible as well.

The obvious prototype of the considered problem is the well-known intractable [1] travelling salesman problem (TSP); see [2–4] et alia. Note papers [5,6] devoted to dynamic programming solutions of the "additive variety" of the TSP. Branch-and-bound methods (see [7]) and various approximate methods (see [4,8] et alia) are widely used.

In real-world applications, there exist problem statements conceptually resembling TSP, however, marked by peculiarities of a specific task; see [2]. Along with the most studied "additive" TSP and its analogs, problems with nonadditive cost aggregation are of interest; see [9,10] et alia. In particular, it is of interest to provide for dynamic programming

solutions of the mentioned "nonadditive" problems with constraints; our paper serves to continue the research in that direction for a more sophisticated bottleneck problem. We study a sequential process made of a finite number of steps; we optimize the cost of the *most costly* step, which leaves us with a minimax problem. In connection with this, let us consider a more substantive instance of the theoretical problem studied below.

Consider the following example of movement cost functions dependent on the list K of pending tasks: for given megalopolises M_k , $k \in K$, assume

exterior movement $c(x, y, K) = \rho(x, y) + a(K) \max_{k \in K} \rho(y, M_k)$,

interior jobs $c_j(\tilde{x}, \tilde{y}, K) = (\rho(\tilde{x}, a_j) + \rho(a_j, \tilde{y}) + b(K) \max_{k \in K \setminus \{j\}} \rho(\tilde{y}, K \setminus \{j\}))$,

where $\rho(\cdot, \cdot)$ is a metric (distance), and $a(K)$ and $b(K)$ are nonnegative functions. These cost functions reflect the demand to construct a route such that in case of an emergency call from some megalopolis that is *not yet traversed* it would be possible to arrive there as quickly as possible *after completing* of the current task. Such circumstances may arise in the operations of a repair brigade that conducts *planned* repairs of the objects (megalopolises) and – in case of an emergency – rescue and reconstruction; the task that is in process at the time of an emergency call is not forfeited because the costs associated with the deployment of the brigade on site (viz., megalopolis) are high, yet comparable with the losses associated with idle time of the object that has suffered an emergency.

1. General Notation and Definitions

We employ the standard set-theoretic notation (quantifiers, propositional connectives, etc.); symbol \triangleq denotes *equality by definition*. Each set, all elements of which are sets themselves, is called a *family*. For every two objects a and b , denote by $\{a; b\}$ the (unique) set that contains a , b , and nothing else. In the case $a = b$, this yields a *singleton* $\{a\} = \{b\}$. If u and v are objects, then $(u, v) \triangleq \{\{u\}; \{u; v\}\}$ [11, p. 67] is the *ordered pair* (OP), the first element of which is u and the second one is v . For an OP z , notation $\text{pr}_1(z)$ denotes its first element and $\text{pr}_2(z)$ denotes its second element; these are uniquely defined by the condition $z = (\text{pr}_1(z), \text{pr}_2(z))$; in case $z \in A \times B$, where A and B are sets, we have $\text{pr}_1(z) \in A$ and $\text{pr}_2(z) \in B$. As usual [12, p. 17], for any three objects a , b , and c , we assume $(a, b, c) \triangleq ((a, b), c)$, which yields the *triple* with the elements a , b , and c . For any three sets A , B , and C , we use the traditional [12, p. 17] convention $A \times B \times C \triangleq (A \times B) \times C$; it obviously means that $(x, y) \in A \times B \times C \ \forall x \in A \ \forall y \in C$. In connection with this, let us also recall the convention we would need later [13, p. 61], which concerns the notation for values of function of three variables: for sets A , B , C , and D , function $h : A \times B \times C \rightarrow D$ and elements $\mu \in A \times B$ and $\nu \in C$, in accordance with the above-mentioned representation of $A \times B \times C$, it is valid to consider the element $h(\mu, \nu) \in D$ to be defined.

As usual, $[0, \infty[\triangleq \{\xi \in \mathbb{R} | 0 \leq \xi\}$ (\mathbb{R} is the real line). For each nonempty set S , denote by $\mathcal{R}_+[S]$ the set of all (nonnegative) functions from S to $[0, \infty[$. As usual, $\mathbb{N} \triangleq \{1; 2; \dots\}$. Assume $\mathbb{N}_0 \triangleq \{0\} \cup \mathbb{N}$ and $\overline{p, q} \triangleq \{i \in \mathbb{N}_0 | (p \leq i) \& (i \leq q)\} \ \forall p \in \mathbb{N}_0 \ \forall q \in \mathbb{N}_0$.

For a nonempty finite set K , let $|K| \in \mathbb{N}$ denote the *power* of the set K ; then, let $(\text{bi})[K]$ denote the set of *all bijections* of the "interval" $\overline{1, |K|}$ onto K . In particular, for a fixed $N \in \mathbb{N}$, let $\mathbb{P} \triangleq (\text{bi})[\overline{1, N}]$ be the set of all *permutations* of the "interval" $\overline{1, N}$; for each $\lambda \in \mathbb{P}$, there exists a permutation $\lambda^{-1} \in \mathbb{P}$ such that $\lambda(\lambda^{-1}(k)) = \lambda^{-1}(\lambda(k)) = k \ \forall k \in \overline{1, N}$.

Denote by $\mathcal{P}(H)$ ($\mathcal{P}'(H)$) the family of all (all nonempty) subsets of set H ; let $\text{Fin}(H)$ be the family of all finite sets from $\mathcal{P}'(H)$.

2. Problem Statement

Here and below, fix a nonempty set X , a point $x^0 \in X$, which is called the *base*, a natural number N , $N \geq 2$, sets $M_1 \in \text{Fin } X, \dots, M_N \in \text{Fin } X$, referred to as *megalopolises*, and relations

$$\mathbb{M}_1 \in \mathcal{P}'(M_1 \times M_1), \dots, \mathbb{M}_N \in \mathcal{P}'(M_N \times M_N). \quad (2.1)$$

For $j \in \overline{1, N}$, OP $z \in \mathbb{M}_j$ describes the possible ways of conducting the interior jobs inside the megalopolis M_j : $\text{pr}_1(z)$ determines the *entry point* and $\text{pr}_2(z)$ determines the *exit point*. We study the issue of organizing of a system of movements

$$\begin{aligned} (x^0) &\rightarrow (\text{pr}_1(z^{(1)}) \in M_{\alpha(1)} \rightsquigarrow \text{pr}_2(z^{(1)}) \in M_{\alpha(1)}) \rightarrow \\ &\rightarrow (\text{pr}_1(z^{(2)}) \in M_{\alpha(2)} \rightsquigarrow \text{pr}_2(z^{(2)}) \in M_{\alpha(2)}) \rightarrow \\ &\rightarrow \dots \rightarrow \\ &\rightarrow (\text{pr}_1(z^{(N)}) \in M_{\alpha(N)} \rightsquigarrow \text{pr}_2(z^{(N)}) \in M_{\alpha(N)}), \end{aligned} \quad (2.2)$$

where α is a permutation of indices from $\overline{1, N}$ and OPs $z^{(1)}, \dots, z^{(N)}$ satisfy the conditions

$$z^{(1)} \in \mathbb{M}_{\alpha(1)}, \dots, z^{(N)} \in \mathbb{M}_{\alpha(N)}. \quad (2.3)$$

In (2.2), we choose the permutation α , in our terms, the *route*, and a tuple $(z^{(1)}, \dots, z^{(N)})$ that agrees with the route in the sense of (2.3); this tuple is called a *track*. Let us stress that the choice of α may be restricted by *precedence constraints*, which would be introduced below. Let us hereinafter assume

$$(x^0 \notin M_j \ \forall j \in \overline{1, N}) \& (M_p \cap M_q = \emptyset \ \forall p \in \overline{1, N} \ \forall q \in \overline{1, N} \setminus \{p\}). \quad (2.4)$$

Conditions (2.4) are relatively common in applications, see [13, Pt. 1,2]. In connection with (2.1), let us make the following conventions:

$$\mathbf{M}_j \triangleq \{\text{pr}_2(z) : z \in \mathbb{M}_j\} \ \forall j \in \overline{1, N}. \quad (2.5)$$

In terms of megalopolises and sets (2.5), let us introduce two nonempty finite subsets of X :

$$\mathbb{X} \triangleq \{x^0\} \cup \left(\bigcup_{i=1}^N M_i \right), \quad \mathbf{X} \triangleq \{x^0\} \cup \left(\bigcup_{i=1}^N \mathbf{M}_i \right); \quad (2.6)$$

clearly, $\mathbf{X} \subset \mathbb{X}$. To define precedence constraints, let us first introduce the set $\mathbf{K} \in \mathcal{P}(\overline{1, N} \times \overline{1, N})$ and call its elements *address pairs*. In an address pair $h \in \mathbf{K}$, the first element $\text{pr}_1(h) \in \overline{1, N}$ is called a *sender*, and the second one $\text{pr}_2(h) \in \overline{1, N}$ is called a *receiver*. The essence of precedence constraints is that for each pair the sender *must* be visited before the receiver. The case $\mathbf{K} = \emptyset$ is not excluded and corresponds to lack of precedence constraints.

Recall that $\mathbb{P} = (\text{bi})[\overline{1, N}]$ is the set of all (complete) routes; it is a nonempty set of cardinality $|\mathbb{P}| = N!$. Clearly [13, Pt. 2],

$$\mathbf{A} \triangleq \left\{ \alpha \in \mathbb{P} \mid \alpha^{-1}(\text{pr}_1(h)) < \alpha^{-1}(\text{pr}_2(h)) \ \forall h \in \mathbf{K} \right\} \quad (2.7)$$

is the set of all routes from \mathbb{P} that are *feasible* in the sense of precedence. Let us assume

$$\forall \mathbf{K}_0 \in \mathcal{P}'(\mathbf{K}) \exists z_0 \in \mathbf{K}_0 : \text{pr}_1(z_0) \neq \text{pr}_2(z) \forall z \in \mathbf{K}_0. \quad (2.8)$$

Assumption (2.8) implies that, in particular, $\text{pr}_1(z) \neq \text{pr}_2(z) \forall z \in \mathbf{K}$. It also implies [13, Pt. 2] that $\mathbf{A} \neq \emptyset$ and, consequently, $\mathbf{A} \in \mathcal{P}'(\mathbb{P})$. Clearly, \mathbf{A} is the set of all routes $\alpha \in \mathbb{P}$ such that

$$\left((\text{pr}_1(z) = \alpha(t_1)) \& (\text{pr}_2(z) = \alpha(t_2)) \right) \Rightarrow (t_1 < t_2)$$

for an address pair $z \in \mathbf{K}$ and "times" $t_1 \in \overline{1, N}$ and $t_2 \in \overline{1, N}$. In addition to the route, we also choose the *track*, or trajectory, which is determined in the sense (2.2) by the OPs $z^{(1)}, \dots, z^{(N)}$, supplemented with the initial OP (x^0, x^0) . To formally define the set of tracks that agree with some route in the sense of (2.2), denote by \mathbb{Z} the set of all tuples $(z_i)_{i \in \overline{0, N}} : \overline{0, N} \rightarrow \mathbb{X} \times \mathbf{X}$. For $\alpha \in \mathbb{P}$, assume

$$\mathcal{Z}_\alpha \triangleq \left\{ (z_i)_{i \in \overline{0, N}} \in \mathbb{Z} \mid (z_0 = (x^0, x^0)) \& (z_t \in \mathbb{M}_{\alpha(t)} \forall t \in \overline{1, N}) \right\}; \quad (2.9)$$

evidently, $\mathcal{Z}_\alpha \in \text{Fin}(\mathbb{Z})$. On the set $\mathbb{X} \times \mathbf{X} \times \mathfrak{N}$, where $\mathfrak{N} \triangleq \mathcal{P}'(\overline{1, N})$ (we call an element of \mathfrak{N} a *task list* or just a *list*), we consider $N + 1$ *cost functions*

$$\mathbf{c} \in \mathcal{R}_+(\mathbb{X} \times \mathbf{X} \times \mathfrak{N}), c_1 \in \mathcal{R}_+(\mathbb{X} \times \mathbf{X} \times \mathfrak{N}), \dots, c_N \in \mathcal{R}_+(\mathbb{X} \times \mathbf{X} \times \mathfrak{N}).$$

For $\alpha \in \mathbb{P}$ and $(z_i)_{i \in \overline{0, N}} \in \mathcal{Z}_\alpha$, assume

$$\begin{aligned} \mathfrak{C}_\alpha[(z_i)_{i \in \overline{0, N}}] \triangleq \max_{t \in \overline{0, N-1}} & \left[\mathbf{c}(\text{pr}_2(z_t), \text{pr}_1(z_{t+1}), \{\alpha(s) : s \in \overline{t+1, N}\}) + \right. \\ & \left. + c_{\alpha(t+1)}(z_{t+1}, \{\alpha(s) : s \in \overline{t+1, N}\}) \right]. \end{aligned} \quad (2.10)$$

We consider (2.10) a *criterion* of the *quality* for a solution $(\alpha, (z_i)_{i \in \overline{0, N}})$ (see (2.2)). Since our choice of the route is restricted by precedence, let *feasible solutions* (FS) be the OPs $(\alpha, (z_i)_{i \in \overline{0, N}})$, $\alpha \in \mathbf{A}, (z_i)_{i \in \overline{0, N}} \in \mathcal{Z}_\alpha$. From the mentioned properties $\mathbf{A} \neq \emptyset$ and $\mathcal{Z}_\alpha \in \text{Fin}(\mathbb{Z})$ for $\alpha \in \mathbf{A}$, we know that FSs form a nonempty finite set. We can now consider the problem

$$\mathfrak{C}_\alpha[(z_i)_{i \in \overline{0, N}}] \rightarrow \min, \alpha \in \mathbf{A}, (z_i)_{i \in \overline{0, N}} \in \mathcal{Z}_\alpha. \quad (2.11)$$

To this problem, we assign its *value* (its extremum)

$$V \triangleq \min_{\alpha \in \mathbf{A}} \min_{(z_i)_{i \in \overline{0, N}} \in \mathcal{Z}_\alpha} \mathfrak{C}_\alpha[(z_i)_{i \in \overline{0, N}}] \in [0, \infty[\quad (2.12)$$

and a nonempty set of optimal FS; an FS $(\alpha^0, (z_i^0)_{i \in \overline{0, N}})$, $\alpha^0 \in \mathbf{A}, (z_i^0)_{i \in \overline{0, N}} \in \mathcal{Z}_{\alpha^0}$, is considered *optimal* for problem (2.11) if $\mathfrak{C}_{\alpha^0}[(z_i^0)_{i \in \overline{0, N}}] = V$.

Let us explore the idea of problem (2.11). Process (2.2) is a number of steps, and at each of them the "system" suffers harmful effects. The latter occur both throughout

exterior movements and *interior jobs* (which take place inside megalopolises). Those steps are

$$x^0 \rightarrow (z^{(1)} \in \mathbb{M}_{\alpha(1)}), \text{pr}_2(z^{(1)}) \rightarrow (z^{(2)} \in \mathbb{M}_{\alpha(2)}), \dots, \text{pr}_2(z^{(N-1)}) \rightarrow (z^{(N)} \in \mathbb{M}_{\alpha(N)}) \quad (2.13)$$

(obviously, (2.13) pertains to the case of $N > 3$). At the end of each step, all harmful effects suffered throughout it are "reset": the system is "decontaminated". It is considered to be important for the doses suffered at each step to not exceed the given threshold connected with the normal operation of the "system".

The objective of movements (2.13) is to "turn off" the sources of harmful effects. In case of movement through radiation fields it may consist of sequential dismantlement of ionizing radiation sources (read disaster cleanup operations similar to those carried out in Chernobyl or Fukushima). Such circumstances may arise when dismantling a decommissioned nuclear power unit if the process is carried out by a special robot that should first of all "survive" to complete the dismantlement of all radiation sources.

3. Extension of the Full Problem and the Dynamic Programming Method

In this section, following [14] and (conceptually) [13, § 4.9], let us consider a natural extension of problem (2.11), which would serve as a base for the necessary version of dynamic programming; the constructions are intentionally abridged. For $K \in \mathfrak{N}$, let $\Sigma[K] \triangleq \{z \in \mathbf{K} \mid (\text{pr}_1(z) \in K) \& (\text{pr}_2(z) \in K)\}$. Based on this, define the mapping \mathbf{I} on \mathfrak{N} by the rule

$$\mathbf{I}(\tilde{K}) \triangleq \tilde{K} \setminus \{\text{pr}_2(z) : z \in \Sigma[\tilde{K}]\} \quad \forall \tilde{K} \in \mathfrak{N}. \quad (3.1)$$

Condition (2.8) directly implies that $\Sigma[\{t\}] = \emptyset$ for $t \in \overline{1, N}$; consequently, $\mathbf{I}(\{t\}) = \{t\}$. In terms of \mathbf{I} (3.1), define the *partial routes* that are feasible in the sense of a *crossing-out* operation connected with \mathbf{I} : For $K \in \mathfrak{N}$, let

$$\begin{aligned} (\mathbf{I} - \text{bi})[K] &\triangleq \left\{ \alpha \in (\text{bi})[K] \mid \alpha(s) \in \mathbf{I}(\{\alpha(t) : t \in \overline{s, |K|}\}) \quad \forall s \in \overline{1, |K|} \right\} = \\ &= \left\{ \alpha \in (\text{bi})[K] \mid \alpha(s) \in \mathbf{I}(K \setminus \{\alpha(t) : t \in \overline{1, s-1}\}) \quad \forall s \in \overline{1, |K|} \right\} = \\ &= \left\{ \alpha \in (\text{bi})[K] \mid (\alpha(1) \in \mathbf{I}(K)) \& (\alpha(s) \in \mathbf{I}(K \setminus \{\alpha(t) : t \in \overline{1, s-1}\})) \quad \forall s \in \overline{2, |K|} \right\}; \end{aligned} \quad (3.2)$$

we assume $\overline{1, 0} = \emptyset$ (for $p \in \mathbb{N}_0$, $q \in \mathbb{N}_0$, and $q < p$, we generally have $\overline{p, q} = \emptyset$; see Section 1). Using properties stated in [13, Pt. 2], in particular, we obtain

$$\begin{aligned} \mathbf{A} &= (\mathbf{I} - \text{bi})[\overline{1, N}] = \\ &= \left\{ \alpha \in \mathbb{P} \mid (\alpha(1) \in \mathbf{I}(\overline{1, N})) \& (\alpha(s) \in \mathbf{I}(\overline{1, N} \setminus \{\alpha(t) : t \in \overline{1, s-1}\})) \quad \forall s \in \overline{2, N} \right\}. \end{aligned} \quad (3.3)$$

It is a special case of (3.2) that corresponds to the main, i.e., "full" problem. Returning to (3.2), let us note (see [13], [14]) that $(\mathbf{I} - \text{bi})[K] \neq \emptyset \quad \forall K \in \mathfrak{N}$. Thus, *crossing-out feasible partial routes* do exist.

Let us now define a *partial track*. For $K \in \mathfrak{N}$, let \mathbb{Z}_K be the set of all tuples $(z_i)_{i \in \overline{0, |K|}} : \overline{0, |K|} \rightarrow \mathbb{X} \times \mathbf{X}$; if, in addition, $x \in \mathbf{X}$ and $\alpha \in (\text{bi})[K]$, let

$$Z(x, K, \alpha) \triangleq \left\{ (z_i)_{i \in \overline{0, |K|}} \in \mathbb{Z}_K \mid (z_0 = (x, x)) \& (z_t \in \mathbb{M}_{\alpha(t)} \forall t \in \overline{1, |K|}) \right\} \in \text{Fin}(\mathbb{Z}_K). \quad (3.4)$$

In view of the definitions just given, define the *partial* (shortened) quality criterions: if $x \in \mathbf{X}$, $K \in \mathfrak{N}$, $\alpha \in (\text{bi})[K]$, and $(z_i)_{i \in \overline{0, |K|}} \in Z(x, K, \alpha)$, let

$$\begin{aligned} \mathfrak{C}_K^{(\alpha)}[(z_i)_{i \in \overline{0, |K|}}] \triangleq \max_{t \in \overline{0, |K|-1}} \left[\mathbf{c} \left(\text{pr}_2(z_t), \text{pr}_1(z_{t+1}), \{ \alpha(s) : s \in \overline{t+1, |K|} \} \right) + \right. \\ \left. + c_{\alpha(t+1)} \left(z_{t+1}, \{ \alpha(s) : s \in \overline{t+1, |K|} \} \right) \right]. \end{aligned} \quad (3.5)$$

In terms of (3.5), we define *partial routing problems* (roughly speaking, *subproblems*): for $x \in \mathbf{X}$ and $K \in \mathfrak{N}$, the problem is

$$\mathfrak{C}_K^{(\alpha)}[(z_i)_{i \in \overline{0, |K|}}] \rightarrow \min, \alpha \in (\mathbf{I} - \text{bi})[K], (z_i)_{i \in \overline{0, |K|}} \in Z(x, K, \alpha). \quad (3.6)$$

Recall that (see, in particular, (3.4)) constraints of each and every problem (3.6) are consistent; each problem has its *value* (its extremum)

$$v(x, K) \triangleq \min_{\alpha \in (\mathbf{I} - \text{bi})[K]} \min_{(z_i)_{i \in \overline{0, |K|}} \in Z(x, K, \alpha)} \mathfrak{C}_K^{(\alpha)}[(z_i)_{i \in \overline{0, |K|}}] \in [0, \infty[\quad (3.7)$$

and a nonempty set of optimal solutions. For $x \in \mathbf{X}$ and $K \in \mathfrak{N}$, FS

$$(\alpha^*, (z_i^*)_{i \in \overline{0, |K|}}), \alpha^* \in \mathbf{A}, (z_i^*)_{i \in \overline{0, |K|}} \in Z(x, K, \alpha^*),$$

is optimal in some problem (3.6) if $\mathfrak{C}_K^{(\alpha^*)}[(z_i^*)_{i \in \overline{0, |K|}}] = v(x, K)$.

Note that in (3.6) and (3.7) we may as well assume $x = x^0$ and $K = \overline{1, N}$. For $\alpha \in \mathbb{P}$ (which means $\alpha \in (\text{bi})[\overline{1, N}]$), by (2.9) and (3.4), we obtain the equality

$$\mathcal{Z}_\alpha = Z(x^0, \overline{1, N}, \alpha) \quad (3.8)$$

(note that $\mathbb{Z} = \mathbb{Z}_{\overline{1, N}}$). Moreover, $\mathbf{A} = (\mathbf{I} - \text{bi})[\overline{1, N}]$ (see [13, Pt. 2]). Then, by (2.12), (3.7), and (3.8), we have the equality

$$V = v(x^0, \overline{1, N}). \quad (3.9)$$

In connection with (3.9), note the obvious consequence of (2.10) and (3.5): for $\alpha \in \mathbb{P}$ and $(z_i)_{i \in \overline{0, N}} \in \mathcal{Z}_\alpha$, we have $\mathfrak{C}_\alpha[(z_i)_{i \in \overline{0, N}}] = \mathfrak{C}_{\overline{1, N}}^{(\alpha)}[(z_i)_{i \in \overline{0, N}}]$, where we also take into account (3.8). Set $v(x, \emptyset) \triangleq 0 \forall x \in \mathbf{X}$. Thus, we have defined the *Bellman function*

$$v : \mathbf{X} \times \mathcal{P}(\overline{1, N}) \rightarrow [0, \infty[. \quad (3.10)$$

Theorem 1. For $x \in \mathbf{X}$ and $K \in \mathfrak{N}$, we have the equality

$$v(x, K) = \min_{j \in \mathbf{I}(K)} \min_{z \in \mathbb{M}_j} \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(z), K) + c_j(z, K); v(\text{pr}_2(z), K \setminus \{j\}) \right\} \right). \quad (3.11)$$

Proof. Fix $x \in \mathbf{X}$ and $K \in \mathfrak{N}$; then, $n \triangleq |K| \in \overline{1, N}$ and $(n = 1) \vee (n \in \overline{2, N})$. These cases are examined separately.

1) In case $n = 1$, we have the equality $K = \{t\}$ for a certain $t \in \overline{1, N}$; consequently, $\mathbf{I}(K) = \{t\}$. Further reasoning is essentially obvious (in the case of $n = 1$); it leads to the conclusion that

$$(n = 1) \Rightarrow \left(v(x, K) = \min_{j \in \mathbf{I}(K)} \min_{z \in \mathbb{M}_j} \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(z), K) + c_j(z, K); v(\text{pr}_2(z), K \setminus \{j\}) \right\} \right) \right). \quad (3.12)$$

2) Let $n \in \overline{2, N}$. Then, $n - 1 \in \overline{1, N - 1}$, $K \setminus \{j\} \in \mathfrak{N}$, and

$$|K \setminus \{j\}| = n - 1 \quad \forall j \in \mathbf{I}(K). \quad (3.13)$$

By (3.7) and (3.13), we have the following equalities:

$$v(\text{pr}_2(z), K \setminus \{j\}) = \min_{\alpha \in (\mathbf{I} - \text{bi})[K \setminus \{j\}]} \min_{(z_i)_{i \in \overline{0, n-1}} \in Z(\text{pr}_2(z), K \setminus \{j\}, \alpha)} \mathfrak{C}_{K \setminus \{j\}}^{(\alpha)}[(z_i)_{i \in \overline{0, n-1}}] \quad \forall j \in \mathbf{I}(K) \quad \forall z \in \mathbb{M}_j. \quad (3.14)$$

This system of equalities is used in the two constructions that follow.

2') In view of (3.7), pick $\alpha^0 \in (\mathbf{I} - \text{bi})[K]$ and $(z_i^0)_{i \in \overline{0, n}} \in Z(x, K, \alpha^0)$ such that

$$v(x, K) = \mathfrak{C}_K^{(\alpha^0)}[(z_i^0)_{i \in \overline{0, n}}]. \quad (3.15)$$

In connection with (3.15), note that

$$\begin{aligned} \mathfrak{C}_K^{(\alpha^0)}[(z_i^0)_{i \in \overline{0, n}}] &= \sup \left(\left\{ \mathbf{c}(x; \text{pr}_1(z_1^0), K) + c_{\alpha^0(1)}(z_1^0, K); \right. \right. \\ &\quad \max_{t \in \overline{1, n-1}} \left[\mathbf{c}(\text{pr}_2(z_t^0), \text{pr}_1(z_{t+1}^0), \{\alpha^0(s) : s \in \overline{t+1, n}\}) + \right. \\ &\quad \left. \left. + c_{\alpha^0(t+1)}(z_{t+1}^0, \{\alpha^0(s) : s \in \overline{t+1, n}\}) \right] \right\} \right). \end{aligned} \quad (3.16)$$

It is of use to remark that $\alpha^0(1) \in \mathbf{I}(K)$ and $z_1^0 \in \mathbb{M}_{\alpha^0(1)}$. Then,

$$\begin{aligned} \min_{j \in \mathbf{I}(K)} \min_{z \in \mathbb{M}_j} \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(z), K) + c_j(z, K); v(\text{pr}_2(z), K \setminus \{j\}) \right\} \right) &\leq \\ &\leq \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(z_1^0), K) + c_{\alpha^0(1)}(z_1^0, K); v(\text{pr}_2(z_1^0), \mathbb{K}) \right\} \right), \end{aligned} \quad (3.17)$$

where $\mathbb{K} \triangleq K \setminus \{\alpha^0(1)\}$. Expression (3.14) now yields the following equality:

$$v(\text{pr}_2(z_1^0), \mathbb{K}) = \min_{\alpha \in (\mathbf{I} - \text{bi})[\mathbb{K}]} \min_{(z_i)_{i \in \overline{0, n-1}} \in Z(\text{pr}_2(z_1^0), \mathbb{K}, \alpha)} \mathfrak{C}_{\mathbb{K}}^{(\alpha)}[(z_i)_{i \in \overline{0, n-1}}]. \quad (3.18)$$

Note that since $|K| = n$, we have $\alpha^0 : \overline{1, n} \twoheadrightarrow K$ (α^0 is surjective), thus $\alpha^0(t+1) \in K$ for $t \in \overline{1, n-1}$. Moreover, since α^0 is injective, $\alpha^0(1) \neq \alpha^0(t) \quad \forall t \in \overline{2, n}$. Then, $\alpha^0(1) \neq$

$\alpha^0(t+1) \forall t \in \overline{1, n-1}$. Finally, $\alpha_0 \triangleq (\alpha^0(t+1))_{t \in \overline{1, n-1}}$ maps $\overline{1, n-1}$ onto \mathbb{K} , i.e., $\alpha_0 : \overline{1, n-1} \rightarrow \mathbb{K}$. Moreover, in view of [15, Proposition 3],

$$\alpha_0 \in (\mathbf{I} - \text{bi})[\mathbb{K}]. \quad (3.19)$$

According to (3.4), $(z_i^0)_{i \in \overline{0, n}} \in \mathbb{Z}_K$ and, in addition,

$$(z_0^0 = (x, x)) \& (z_t^0 \in \mathbb{M}_{\alpha^0(t)} \forall t \in \overline{1, n}). \quad (3.20)$$

Then, $\text{pr}_2(z_1^0) \in \mathbf{X}$ and hence, in view of (3.4) and (3.19) and the equality $|\mathbb{K}| = n-1$, we have

$$\begin{aligned} Z(\text{pr}_2(z_1^0), \mathbb{K}, \alpha_0) = & \left\{ (z_i)_{i \in \overline{0, n-1}} \in \mathbb{Z}_{\mathbb{K}} \mid (z_0 = (\text{pr}_2(z_1^0), \text{pr}_2(z_1^0))) \& \right. \\ & \left. \& (z_t \in \mathbb{M}_{\alpha_0(t)} \forall t \in \overline{1, n-1}) \right\}. \end{aligned} \quad (3.21)$$

Consider a tuple $(\hat{z}_t^0)_{t \in \overline{0, n-1}} : \overline{0, n-1} \rightarrow \mathbb{X} \times \mathbf{X}$ defined in the following way:

$$\left(\hat{z}_0^0 = (\text{pr}_2(z_1^0), \text{pr}_2(z_1^0)) \right) \& \left(\hat{z}_t^0 \triangleq z_{t+1}^0 \forall t \in \overline{1, n-1} \right). \quad (3.22)$$

Since $|\mathbb{K}| = n-1$, taking into account (3.21) and (3.22), we obtain

$$(\hat{z}_t^0)_{t \in \overline{0, n-1}} \in Z(\text{pr}_2(z_1^0), \mathbb{K}, \alpha_0). \quad (3.23)$$

From (3.18), (3.19), and (3.23), it follows that

$$\begin{aligned} v(\text{pr}_2(z_1^0), \mathbb{K}) \leq & \mathfrak{C}_{\mathbb{K}}^{(\alpha_0)}[(\hat{z}_i^0)_{i \in \overline{0, n-1}}] = \max_{t \in \overline{0, n-2}} \left[\mathbf{c}(\text{pr}_2(\hat{z}_t^0), \text{pr}_1(\hat{z}_{t+1}^0), \{\alpha_0(s) : \right. \\ & \left. s \in \overline{t+1, n-1}\}) + c_{\alpha_0(t+1)}(\hat{z}_{t+1}^0, \{\alpha_0(s) : s \in \overline{t+1, n-1}\}) \right]. \end{aligned} \quad (3.24)$$

Since $\{\alpha_0(s) : s \in \overline{t+1, n-1}\} = \{\alpha^0(s+1) : s \in \overline{t+1, n-1}\} = \{\alpha^0(l) : l \in \overline{t+2, n}\}$, the expression in the right-hand side of (3.24) takes the form

$$\begin{aligned} \mathfrak{C}_{\mathbb{K}}^{(\alpha_0)}[(\hat{z}_t^0)_{t \in \overline{0, n-1}}] = & \max_{t \in \overline{0, n-2}} \left[\mathbf{c}(\text{pr}_2(\hat{z}_t^0), \text{pr}_1(\hat{z}_{t+1}^0), \{\alpha^0(l) : l \in \overline{t+2, n}\}) + \right. \\ & \left. + c_{\alpha^0(t+2)}(\hat{z}_{t+1}^0, \{\alpha^0(l) : l \in \overline{t+2, n}\}) \right]. \end{aligned} \quad (3.25)$$

Expression (3.25) implies the following obvious property:

$$\begin{aligned} \mathfrak{C}_{\mathbb{K}}^{(\alpha_0)}[(\hat{z}_t^0)_{t \in \overline{0, n-1}}] = & \max_{t \in \overline{0, n-2}} \left[\mathbf{c}(\text{pr}_2(z_{t+1}^0), \text{pr}_1(z_{t+2}^0), \{\alpha^0(l) : l \in \overline{t+2, n}\}) + \right. \\ & \left. + c_{\alpha^0(t+2)}(z_{t+2}^0, \{\alpha^0(l) : l \in \overline{t+2, n}\}) \right] = \max_{\theta \in \overline{1, n-1}} \left[\mathbf{c}(\text{pr}_2(z_{\theta}^0), \text{pr}_1(z_{\theta+1}^0), \{\alpha^0(l) : \right. \\ & \left. l \in \overline{\theta+1, n}\}) + c_{\alpha^0(\theta+1)}(z_{\theta+1}^0, \{\alpha^0(l) : l \in \overline{\theta+1, n}\}) \right]. \end{aligned} \quad (3.26)$$

From (3.24) and (3.26), we obtain the estimate

$$\begin{aligned} v(\text{pr}_2(z_1^0), \mathbb{K}) \leq & \max_{\theta \in \overline{1, n-1}} \left[\mathbf{c}(\text{pr}_2(z_{\theta}^0), \text{pr}_1(z_{\theta+1}^0), \{\alpha^0(l) : l \in \overline{\theta+1, n}\}) + \right. \\ & \left. + c_{\alpha^0(\theta+1)}(z_{\theta+1}^0, \{\alpha^0(l) : l \in \overline{\theta+1, n}\}) \right]. \end{aligned} \quad (3.27)$$

From (3.17) and (3.27), we have

$$\begin{aligned} & \min_{j \in \mathbf{I}(K)} \min_{z \in \mathbb{M}_j} \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(z), K) + c_j(z, K); v(\text{pr}_2(z), K \setminus \{j\}) \right\} \right) \leq \\ & \leq \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(z_1^0), K) + c_{\alpha^0(1)}(z_1^0, K); \max_{t \in \overline{1, n-1}} \left[\mathbf{c}(\text{pr}_2(z_t^0), \text{pr}_1(z_{t+1}^0), \right. \right. \right. \\ & \left. \left. \left. \{\alpha^0(l) : l \in \overline{t+1, n}\} \right) + c_{\alpha^0(t+1)}(z_{t+1}^0, \{\alpha^0(l) : l \in \overline{t+1, n}\}) \right] \right\} \right). \end{aligned} \quad (3.28)$$

Expressions (3.16) and (3.28) yield the inequality

$$\min_{j \in \mathbf{I}(K)} \min_{z \in \mathbb{M}_j} \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(z), K) + c_j(z, K); v(\text{pr}_2(z), K \setminus \{j\}) \right\} \right) \leq \mathfrak{C}_K^{(\alpha^0)}[(z_t^0)_{t \in \overline{0, n}}]. \quad (3.29)$$

From (3.15) and (3.29), we obtain the estimate

$$\min_{j \in \mathbf{I}(K)} \min_{z \in \mathbb{M}_j} \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(z), K) + c_j(z, K); v(\text{pr}_2(z), K \setminus \{j\}) \right\} \right) \leq v(x, K). \quad (3.30)$$

Choose and fix an index q and its OP of entry and exit points \mathbf{z}

$$q \in \mathbf{I}(K), \quad (3.31)$$

$$\mathbf{z} \in \mathbb{M}_q, \quad (3.32)$$

such that

$$\begin{aligned} & \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(\mathbf{z}), K) + c_q(\mathbf{z}, K); v(\text{pr}_2(\mathbf{z}), K \setminus \{q\}) \right\} \right) = \\ & = \min_{j \in \mathbf{I}(K)} \min_{z \in \mathbb{M}_j} \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(z), K) + c_j(z, K); v(\text{pr}_2(z), K \setminus \{j\}) \right\} \right). \end{aligned} \quad (3.33)$$

For brevity, set

$$Q \triangleq K \setminus \{q\}. \quad (3.34)$$

Expression (3.31) implies that, in particular, $q \in K$, whence (see (3.34)) $|Q| = n - 1 \geq 1$. From (3.33) and (3.34) it follows that

$$\begin{aligned} & \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(\mathbf{z}), K) + c_q(\mathbf{z}, K); v(\text{pr}_2(\mathbf{z}), Q) \right\} \right) = \\ & = \min_{j \in \mathbf{I}(K)} \min_{z \in \mathbb{M}_j} \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(z), K) + c_j(z, K); v(\text{pr}_2(z), K \setminus \{j\}) \right\} \right). \end{aligned} \quad (3.35)$$

From (3.32), we now obtain

$$\text{pr}_2(\mathbf{z}) \in \mathbf{M}_q \quad (3.36)$$

and, in particular, $\text{pr}_2(\mathbf{z}) \in \mathbf{X}$; see (2.6), (3.36). Then,

$$(\text{pr}_2(\mathbf{z}), Q) \in \mathbf{X} \times \mathfrak{N}. \quad (3.37)$$

Thus, according to (3.7) and (3.37),

$$v(\text{pr}_2(\mathbf{z}), Q) = \min_{\alpha \in (\mathbf{I} - \text{bi})[Q]} \min_{(z_i)_{i \in \overline{0, n-1}} \in Z(\text{pr}_2(\mathbf{z}), Q, \alpha)} \mathfrak{C}_Q^{(\alpha)}[(z_i)_{i \in \overline{0, n-1}}]; \quad (3.38)$$

here, in view of (3.5), we have the equalities

$$\begin{aligned} \mathfrak{C}_Q^{(\alpha)} \left[(z_i)_{i \in \overline{0, n-1}} \right] &= \max_{t \in \overline{0, n-2}} \left[\mathbf{c}(\text{pr}_2(z_t), \text{pr}_1(z_{t+1}), \{\alpha(s) : s \in \overline{t+1, n-1}\}) + \right. \\ &\quad \left. + c_{\alpha(t+1)}(z_{t+1}, \{\alpha(s) : s \in \overline{t+1, n-1}\}) \right] \\ \forall \alpha \in (\mathbf{I} - \text{bi})[Q] \forall (z_i)_{i \in \overline{0, n-1}} \in Z(\text{pr}_2(\mathbf{z}), Q, \alpha). \end{aligned} \quad (3.39)$$

In accordance with (3.38), let us choose and fix a route β and a track $(h_i)_{i \in \overline{0, n-1}}$

$$\beta \in (\mathbf{I} - \text{bi})[Q], \quad (3.40)$$

$$(h_i)_{i \in \overline{0, n-1}} \in Z(\text{pr}_2(\mathbf{z}), Q, \beta), \quad (3.41)$$

such that

$$v(\text{pr}_2(\mathbf{z}), Q) = \mathfrak{C}_Q^{(\beta)} \left[(h_i)_{i \in \overline{0, n-1}} \right], \quad (3.42)$$

where (see (3.39), (3.40), (3.41))

$$\begin{aligned} \mathfrak{C}_Q^{(\beta)} \left[(h_i)_{i \in \overline{0, n-1}} \right] &= \max_{t \in \overline{0, n-2}} \left[\mathbf{c}(\text{pr}_2(h_t), \text{pr}_1(h_{t+1}), \{\beta(s) : s \in \overline{t+1, n-1}\}) + \right. \\ &\quad \left. + c_{\beta(t+1)}(h_{t+1}, \{\beta(s) : s \in \overline{t+1, n-1}\}) \right]. \end{aligned} \quad (3.43)$$

From (3.42) and (3.43), we obtain the following equality

$$\begin{aligned} v(\text{pr}_2(\mathbf{z}), Q) &= \max_{t \in \overline{0, n-2}} \left[\mathbf{c}(\text{pr}_2(h_t), \text{pr}_1(h_{t+1}), \{\beta(s) : s \in \overline{t+1, n-1}\}) + \right. \\ &\quad \left. + c_{\beta(t+1)}(h_{t+1}, \{\beta(s) : s \in \overline{t+1, n-1}\}) \right]. \end{aligned} \quad (3.44)$$

Expressions (3.33), (3.35), and (3.44) imply that

$$\begin{aligned} \min_{j \in \mathbf{I}(K)} \min_{z \in \mathbb{M}_j} \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(z), K) + c_j(z, K); v(\text{pr}_2(z), K \setminus \{j\}) \right\} \right) &= \\ = \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(\mathbf{z}), K) + c_q(\mathbf{z}, K); \max_{t \in \overline{0, n-2}} \left[\mathbf{c}(\text{pr}_2(h_t), \text{pr}_1(h_{t+1}), \right. \right. \right. & \\ \left. \left. \left. \{\beta(s) : s \in \overline{t+1, n-1}\}) + c_{\beta(t+1)}(h_{t+1}, \{\beta(s) : s \in \overline{t+1, n-1}\}) \right] \right\} \right). \end{aligned} \quad (3.45)$$

Note that $q \in K$ and, moreover (see (3.34), (3.40)), $\beta(j) \in K \forall j \in \overline{1, n-1}$. In view of that, let us introduce a mapping

$$\hat{\beta} : \overline{1, n} \rightarrow K, \quad (3.46)$$

defined by the following conditions:

$$(\hat{\beta}(1) \triangleq q) \& (\hat{\beta}(j) \triangleq \beta(j-1) \forall j \in \overline{2, n}). \quad (3.47)$$

Let us show that $\hat{\beta} \in (\text{bi})[K]$. To this end, let us start by noticing that, since β maps $\overline{1, n-1}$ onto Q , we have $\forall j \in Q \exists s \in \overline{2, n} : j = \hat{\beta}(j)$. In view of (3.34) and (3.47), we obtain $K \subset \{\hat{\beta}(l) : l \in \overline{1, n}\}$, i.e., $K = \{\hat{\beta}(l) : l \in \overline{1, n}\}$ (we also took into account (3.46)).

Thus, $\hat{\beta}$ is (see (3.46)) a surjection of $\overline{1, n}$ onto K . Let us check the injectivity of $\hat{\beta}$. Fix indices $j_* \in \overline{1, n}$ and $j^* \in \overline{1, n}$ such that

$$\hat{\beta}(j_*) = \hat{\beta}(j^*). \quad (3.48)$$

Let $\nu \triangleq \hat{\beta}(j_*)$. Then, $\nu \in K$ and $\hat{\beta}(j^*) = \nu$. The two cases below are mutually exclusive:

$$(\nu = q) \vee (\nu \in Q). \quad (3.49)$$

a) Let $\nu = q$. Since in view of (3.37), (3.40), and (3.47) we have $\hat{\beta}(j) \in Q$ for $j \in \overline{2, n}$, we now see that $\hat{\beta}(l) \neq q \forall l \in \overline{2, n}$. As $\hat{\beta}(j_*) = \hat{\beta}(j^*) = q$ in the considered case, we obtain $j_* = 1 = j^*$. We proved that

$$(\nu = q) \Rightarrow (j_* = j^*). \quad (3.50)$$

b) Let $\nu \in Q$. Note that in this case $(\hat{\beta}(j_*) \in Q) \& (\hat{\beta}(j^*) \in Q)$. We can therefore conclude from (3.34) and (3.47) that $j_* \neq 1$ and $j^* \neq 1$. Then, $(j_* \in \overline{2, n}) \& (j^* \in \overline{2, n})$ and hence (see (3.48)) $\beta(j_* - 1) = \hat{\beta}(j_*) = \hat{\beta}(j^*) = \beta(j^* - 1)$; since β is injective (see (3.40)), $j_* = j^*$. Thus, $(\nu \in Q) \Rightarrow (j_* = j^*)$. Taking into account (3.49) and (3.50), we obtain the equality $j_* = j^*$ for all cases. Thus (see (3.48)), $(\hat{\beta}(j_*) = \hat{\beta}(j^*)) \Rightarrow (j_* = j^*)$. Since the choice of j_* and j^* was arbitrary, we have proved the injectivity of $\hat{\beta}$, whence, finally,

$$\hat{\beta} \in (\text{bi})[K], \quad (3.51)$$

i.e., $\hat{\beta}$ is a partial route. Let us prove that it is crossing-out feasible:

$$\hat{\beta} \in (\mathbf{I} - \text{bi})[K]. \quad (3.52)$$

Indeed, let $\gamma \in \overline{1, n}$ (recall that $|K| = n$); consider the set $\Gamma \triangleq \{\hat{\beta}(t) : t \in \overline{\gamma, n}\}$. The two following cases are mutually exclusive:

$$(\gamma = 1) \vee (\gamma \in \overline{2, n}). \quad (3.53)$$

a') Firstly, let $\gamma = 1$. Then, $\overline{\gamma, n} = \overline{1, n}$ and, by virtue of surjectivity of $\hat{\beta}$, we have $\Gamma = \{\hat{\beta}(t) : t \in \overline{1, n}\} = K$; hence, in view of (3.31), $q \in \mathbf{I}(\Gamma)$ and, according to (3.47), $\hat{\beta}(1) \in \mathbf{I}(\Gamma)$, where $\hat{\beta}(1) = \hat{\beta}(\gamma)$. Thus, we have proved the implication

$$(\gamma = 1) \Rightarrow (\hat{\beta}(\gamma) \in \mathbf{I}(\Gamma)). \quad (3.54)$$

b') Now, let $\gamma \in \overline{2, n}$. Then, $\gamma - 1 \in \overline{1, n - 1}$. Therefore, we can consider the index

$$\hat{\beta}(\gamma) = \beta(\gamma - 1) \in Q \quad (3.55)$$

(see (3.40)). Then, according to (3.2), (3.40), and (3.55),

$$\hat{\beta}(\gamma) \in \mathbf{I}(\{\beta(t) : t \in \overline{\gamma - 1, n - 1}\}). \quad (3.56)$$

Let us introduce the set

$$\Gamma_0 \triangleq \{\beta(t) : t \in \overline{\gamma - 1, n - 1}\} \in \mathcal{P}'(Q). \quad (3.57)$$

Show that $\Gamma = \Gamma_0$. Indeed, let $y_* \in \Gamma$ and let $t_* \in \overline{\gamma, n}$ be such that

$$y_* = \hat{\beta}(t_*). \quad (3.58)$$

Then, since $2 \leq \gamma$, we also know that $t_* \in \overline{2, n}$, $t_* - 1 \in \overline{1, n-1}$, and $\hat{\beta}(t_*) = \beta(t_* - 1)$, i.e., $y_* = \beta(t_* - 1)$ in view of (3.58). On the other hand, we chose t_* such that $t_* - 1 \in \overline{\gamma - 1, n - 1}$, whence (see (3.57)) $\beta(t_* - 1) \in \Gamma_0$ and, therefore, $y_* \in \Gamma_0$. Since the choice of y_* was arbitrary, we have found out that

$$\Gamma \subset \Gamma_0. \quad (3.59)$$

Now, let $y_0 \in \Gamma_0$ and let $t_0 \in \overline{\gamma - 1, n - 1}$ implement the equality

$$y_0 = \beta(t_0). \quad (3.60)$$

In addition, $t_0 + 1 \in \overline{\gamma, n}$ and, by definition of Γ , we have the inclusion

$$\hat{\beta}(t_0 + 1) \in \Gamma. \quad (3.61)$$

In addition to this, $2 \leq \gamma \leq t_0 + 1$ and, therefore, $t_0 + 1 \in \overline{2, n}$, hence (see (3.47)) $\hat{\beta}(t_0 + 1) = \beta(t_0)$ and, by (3.60), $y_0 = \hat{\beta}(t_0 + 1)$; therefore, (3.61) implies that $y_0 \in \Gamma$. Since the choice of y_0 was arbitrary, we have proven that $\Gamma_0 \subset \Gamma$, and hence, in view of (3.59), we obtain the desired

$$\Gamma = \Gamma_0, \quad (3.62)$$

where, in accordance with (3.56) and (3.57), we have the inclusion $\hat{\beta}(\gamma) \in \mathbf{I}(\Gamma_0)$, which means (see (3.62)) that $\hat{\beta}(\gamma) \in \mathbf{I}(\Gamma)$ for $\gamma \in \overline{2, n}$ as well. Thus, we have proven the implication $(\gamma \in \overline{2, n}) \Rightarrow (\hat{\beta}(\gamma) \in \mathbf{I}(\Gamma))$. In view of (3.53) and (3.54), we have the inclusion $\hat{\beta}(\gamma) \in \mathbf{I}(\Gamma)$ in all possible cases, i.e., $\hat{\beta}(\gamma) \in \mathbf{I}(\{\hat{\beta}(t) : t \in \overline{\gamma, n}\})$. Since the choice of γ was arbitrary, we have proven that

$$\hat{\beta}(s) \in \mathbf{I}(\{\hat{\beta}(t) : t \in \overline{s, n}\}) \quad \forall s \in \overline{1, n}. \quad (3.63)$$

From (3.51) and (3.63), we obtain (see (3.2)) the desired property (3.52), i.e., now $\hat{\beta} \in (\mathbf{I} - \text{bi})[K]$. Note that (3.4) and (3.41) imply that $(h_i)_{i \in \overline{0, n-1}} \in \mathbb{Z}_Q$, i.e.,

$$(h_i)_{i \in \overline{0, n-1}} : \overline{0, n-1} \rightarrow \mathbb{X} \times \mathbf{X}. \quad (3.64)$$

According to (3.4) and (3.41),

$$\left(h_0 = (\text{pr}_2(\mathbf{z}), \text{pr}_2(\mathbf{z})) \right) \& \left(h_t \in \mathbb{M}_{\beta(t)} \quad \forall t \in \overline{1, n-1} \right). \quad (3.65)$$

Note that from (3.65) we have the OPs

$$h_{t-1} \in \mathbb{M}_{\beta(t-1)} \quad \forall t \in \overline{2, n}. \quad (3.66)$$

We also have (see (3.47), (3.66)) the following property:

$$h_{t-1} \in \mathbb{M}_{\hat{\beta}(t)} \quad \forall t \in \overline{2, n}. \quad (3.67)$$

Note that $\overline{1, n} = \{1\} \cup \overline{2, n}$ and $\overline{0, n} = \{0\} \cup \overline{1, n} = \{0\} \cup \{1\} \cup \overline{2, n}$. In view of (3.64), we may now consider a tuple $(\hat{h}_t)_{t \in \overline{0, n}} : \overline{0, n} \rightarrow \mathbb{X} \times \mathbf{X}$ defined by conditions

$$(\hat{h}_0 \triangleq (x, x)) \& (\hat{h}_1 \triangleq \mathbf{z}) \& (\hat{h}_t \triangleq h_{t-1} \ \forall t \in \overline{2, n}). \quad (3.68)$$

Clearly (since $n = |K|$),

$$(\hat{h}_t)_{t \in \overline{0, n}} \in \mathbb{Z}_K. \quad (3.69)$$

Then, (3.32) and (3.68) imply that $\hat{h}_1 \in \mathbb{M}_q$, whence, taking into account (3.47), we obtain

$$\hat{h}_1 \in \mathbb{M}_{\hat{\beta}(1)}. \quad (3.70)$$

Then, (3.4) and (3.41) imply that $h_t \in \mathbb{M}_{\beta(t)} \ \forall t \in \overline{1, n-1}$. Then, expression (3.68), yields $\hat{h}_t \in \mathbb{M}_{\beta(t-1)} \ \forall t \in \overline{2, n}$. Using (3.47), we conclude that $\hat{h}_t \in \mathbb{M}_{\hat{\beta}(t)} \ \forall t \in \overline{2, n}$. In view of (3.70) and the last property, we obtain the system of inclusions

$$\hat{h}_t \in \mathbb{M}_{\hat{\beta}(t)} \ \forall t \in \overline{1, n}. \quad (3.71)$$

From (3.68), (3.69), and (3.71), we conclude that

$$(\hat{h}_t)_{t \in \overline{0, n}} \in \mathbb{Z}_K : (\hat{h}_0 = (x, x)) \& (\hat{h}_t \in \mathbb{M}_{\hat{\beta}(t)} \ \forall t \in \overline{1, n}). \quad (3.72)$$

But in this case, (3.4) and (3.72) imply that

$$(\hat{h}_t)_{t \in \overline{0, n}} \in Z(x, K, \hat{\beta}). \quad (3.73)$$

In accordance with (3.52) and (3.73), we see that the OP $(\hat{\beta}, (\hat{h}_t)_{t \in \overline{0, n}})$ is an FS of the considered *partial problem*, therefore (see (3.7)),

$$v(x, K) \leq \mathfrak{C}_K^{(\hat{\beta})} \left[(\hat{h}_t)_{t \in \overline{0, n}} \right]. \quad (3.74)$$

In view of (3.5), (3.52), and (3.73), we have

$$\begin{aligned} \mathfrak{C}_K^{(\hat{\beta})} \left[(\hat{h}_t)_{t \in \overline{0, n}} \right] &= \max_{t \in \overline{0, n-1}} \left[\mathbf{c}(\text{pr}_2(\hat{h}_t), \text{pr}_1(\hat{h}_{t+1}), \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) \right. \\ &\quad \left. + c_{\hat{\beta}(t+1)}(\hat{h}_{t+1}, \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) \right] \in [0, \infty[. \end{aligned} \quad (3.75)$$

Consequently, from (3.68) and (3.75), we conclude that

$$\begin{aligned} \mathfrak{C}_K^{(\hat{\beta})} \left[(\hat{h}_t)_{t \in \overline{0, n}} \right] &= \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(\hat{h}_1), \{\hat{\beta}(s) : s \in \overline{1, n}\}) + c_{\hat{\beta}(1)}(\hat{h}_1, \{\hat{\beta}(s) : s \in \overline{1, n}\}); \right. \right. \\ &\quad \left. \left. \max_{t \in \overline{1, n-1}} \left[\mathbf{c}(\text{pr}_2(\hat{h}_t), \text{pr}_1(\hat{h}_{t+1}), \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) + c_{\hat{\beta}(t+1)}(\hat{h}_{t+1}, \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) \right] \right\} \right). \end{aligned}$$

In view of $n = |K|$, expression (3.51) yields $\{\hat{\beta}(s) : s \in \overline{1, n}\} = K$. Using that, from (3.46), (3.68), (3.74), and (3.75), we obtain the inequality

$$\begin{aligned} v(x, K) &\leq \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(\mathbf{z}), K) + c_q(\mathbf{z}, K); \max_{t \in \overline{1, n-1}} \left[\mathbf{c}(\text{pr}_2(\hat{h}_t), \text{pr}_1(\hat{h}_{t+1}), \right. \right. \right. \\ &\quad \left. \left. \left. \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) + c_{\hat{\beta}(t+1)}(\hat{h}_{t+1}, \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) \right] \right\} \right). \end{aligned} \quad (3.76)$$

Consider the transformation of the expression

$$\max_{t \in \overline{1, n-1}} \left[\mathbf{c}(\text{pr}_2(\hat{h}_t), \text{pr}_1(\hat{h}_{t+1}), \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) + c_{\hat{\beta}(t+1)}(\hat{h}_{t+1}, \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) \right] \in [0, \infty[.$$

To this end, note that, according to (3.68),

$$\mathbf{c}(\text{pr}_2(\hat{h}_1), \text{pr}_1(\hat{h}_2), \{\hat{\beta}(s) : s \in \overline{2, n}\}) = \mathbf{c}(\text{pr}_2(\mathbf{z}), \text{pr}_1(h_1), \{\hat{\beta}(s) : s \in \overline{2, n}\}). \quad (3.77)$$

However, (3.65) implies that $\text{pr}_2(\mathbf{z}) = \text{pr}_2(h_0)$. Then, (3.77) yields the equality

$$\mathbf{c}(\text{pr}_2(\hat{h}_1), \text{pr}_1(\hat{h}_2), \{\hat{\beta}(s) : s \in \overline{2, n}\}) = \mathbf{c}(\text{pr}_2(h_0), \text{pr}_1(h_1), \{\hat{\beta}(s) : s \in \overline{2, n}\}). \quad (3.78)$$

Moreover, (3.65) implies that (see (3.47)) $c_{\hat{\beta}(2)}(\hat{h}_2, \{\hat{\beta}(s) : s \in \overline{2, n}\}) = c_{\beta(1)}(h_1, \{\hat{\beta}(s) : s \in \overline{2, n}\})$. Taking into account (3.78), we obtain

$$\begin{aligned} & \mathbf{c}(\text{pr}_2(\hat{h}_1), \text{pr}_1(\hat{h}_2), \{\hat{\beta}(s) : s \in \overline{2, n}\}) + c_{\hat{\beta}(2)}(\hat{h}_2, \{\hat{\beta}(s) : s \in \overline{2, n}\}) = \\ & = \mathbf{c}(\text{pr}_2(h_0), \text{pr}_1(h_1), \{\hat{\beta}(s) : s \in \overline{2, n}\}) + c_{\beta(1)}(h_1, \{\hat{\beta}(s) : s \in \overline{2, n}\}). \end{aligned}$$

This equality can be transformed into the following expression: if $t \in \overline{1, n-1}$, then, for $t = 1$,

$$\begin{aligned} & \mathbf{c}(\text{pr}_2(\hat{h}_t), \text{pr}_1(\hat{h}_{t+1}), \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) + c_{\hat{\beta}(t+1)}(\hat{h}_{t+1}, \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) = \\ & = \mathbf{c}(\text{pr}_2(h_{t-1}), \text{pr}_1(h_t), \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) + c_{\beta(t)}(h_t, \{\hat{\beta}(s) : s \in \overline{t+1, n}\}). \end{aligned} \quad (3.79)$$

Choose arbitrary $\tau \in \overline{2, n-1}$. Then, $\tau - 1 \in \overline{1, n-2}$ and $\tau + 1 \in \overline{3, n}$. From (3.47), we now have, in particular,

$$\hat{\beta}(\tau + 1) = \beta(\tau) \in K. \quad (3.80)$$

Then, (3.68) implies the following representations:

$$(\hat{h}_\tau = h_{\tau-1}) \& (\hat{h}_{\tau+1} = h_\tau). \quad (3.81)$$

Thus (see (3.80), (3.81)) we obtain

$$\begin{aligned} & \mathbf{c}(\text{pr}_2(\hat{h}_\tau), \text{pr}_1(\hat{h}_{\tau+1}), \{\hat{\beta}(s) : s \in \overline{\tau+1, n}\}) + c_{\hat{\beta}(\tau+1)}(\hat{h}_{\tau+1}, \{\hat{\beta}(s) : s \in \overline{\tau+1, n}\}) = \\ & = \mathbf{c}(\text{pr}_2(h_{\tau-1}), \text{pr}_1(h_\tau), \{\hat{\beta}(s) : s \in \overline{\tau+1, n}\}) + c_{\beta(\tau)}(h_\tau, \{\hat{\beta}(s) : s \in \overline{\tau+1, n}\}). \end{aligned} \quad (3.82)$$

Since the choice of τ was arbitrary, we proved that

$$\begin{aligned} & \mathbf{c}(\text{pr}_2(\hat{h}_t), \text{pr}_1(\hat{h}_{t+1}), \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) + c_{\hat{\beta}(t+1)}(\hat{h}_{t+1}, \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) \\ & = \mathbf{c}(\text{pr}_2(h_{t-1}), \text{pr}_1(h_t), \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) + \\ & + c_{\beta(t)}(h_t, \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) \quad \forall t \in \overline{2, n-1}. \end{aligned} \quad (3.83)$$

From (3.79) and (3.83), we conclude that

$$\begin{aligned} & \mathbf{c}(\text{pr}_2(\hat{h}_t), \text{pr}_1(\hat{h}_{t+1}), \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) + c_{\hat{\beta}(t+1)}(\hat{h}_{t+1}, \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) = \\ & = \mathbf{c}(\text{pr}_2(h_{t-1}), \text{pr}_1(h_t), \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) + \\ & + c_{\beta(t)}(h_t, \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) \quad \forall t \in \overline{1, n-1}. \end{aligned} \quad (3.84)$$

Let us now consider the set $\{\hat{\beta}(s) : s \in \overline{t^\sharp + 1, n}\}$ for $t^\sharp \in \overline{1, n-1}$. Then, for $s \in \overline{t^\sharp + 1, n}$, we have $s \in \overline{2, n}$, whence (see (3.47)) $\hat{\beta}(s) = \beta(s-1)$, where $s-1 \in \overline{t^\sharp, n-1}$. In particular, we have $s-1 \in \overline{1, n-1} \forall s \in \overline{t^\sharp + 1, n}$. Thus (see (3.40)) we may define the indices

$$\beta(s-1) \in Q \forall s \in \overline{t^\sharp + 1, n}. \quad (3.85)$$

Therefore, we have the set $\{\beta(s-1) : s \in \overline{t^\sharp + 1, n}\} \in \mathcal{P}'(Q)$; recall that $t^\sharp + 1 \leq n$. Returning to (3.85), note that $\overline{t^\sharp, n-1}$ is a nonempty subset of $\overline{1, n-1}$ (by the choice of t^\sharp), thus $\beta(l) \in Q \forall l \in \overline{t^\sharp, n-1}$. Therefore, we also have the set $\mathbb{B} \triangleq \{\beta(l) : l \in \overline{t^\sharp, n-1}\} \in \mathcal{P}'(Q)$ and the set $\mathbf{B} \triangleq \{\hat{\beta}(l) : l \in \overline{t^\sharp + 1, n}\} \in \mathcal{P}'(K)$. In any case, \mathbb{B} and \mathbf{B} are not empty. Let us prove that they are identical. Let $p \in \mathbb{B}$. Then, $p \in Q$ and, for a certain $l_* \in \overline{t^\sharp, n-1}$,

$$p = \beta(l_*). \quad (3.86)$$

Then, $l_* \in \overline{1, n-1}$, $l^* \triangleq l_* + 1 \in \overline{2, n}$ and, according to (3.46), $\hat{\beta}(l^*) \in K$; however, (3.47) implies that $\hat{\beta}(l^*) = \beta(l^* - 1) = \beta(l_*) = p$. Nevertheless, $l^* \in \overline{t^\sharp + 1, n}$, whence $\hat{\beta}(l^*) \in \mathbf{B}$. Thus $p \in \mathbf{B}$, which completes the testing of the inclusion

$$\mathbb{B} \subset \mathbf{B}. \quad (3.87)$$

Now, let $\tilde{p} \in \mathbf{B}$. Then, $\tilde{p} \in K$ and, for a certain $\tilde{l}_* \in \overline{t^\sharp + 1, n}$, we have the equality

$$\tilde{p} = \hat{\beta}(\tilde{l}_*). \quad (3.88)$$

In particular, $\tilde{l}_* \in \overline{2, n}$, $\tilde{l}^* \triangleq \tilde{l}_* - 1 \in \overline{1, n-1}$ and, according to (3.47), $\tilde{p} = \hat{\beta}(\tilde{l}_*) = \beta(\tilde{l}_* - 1) = \beta(\tilde{l}^*)$. However, by the choice of \tilde{l}_* , we have $\tilde{l}^* = \tilde{l}_* - 1 \in \overline{t^\sharp, n-1}$, whence (see (3.88)) $\tilde{p} = \beta(\tilde{l}^*) \in \mathbb{B}$. We established the inclusion

$$\mathbf{B} \subset \mathbb{B}. \quad (3.89)$$

In view of (3.87) and (3.89), we obtain the desired equality

$$\mathbb{B} = \mathbf{B}. \quad (3.90)$$

In other words, (3.90) implies that $\{\hat{\beta}(s) : s \in \overline{t^\sharp + 1, n}\} = \{\beta(l) : l \in \overline{t^\sharp, n-1}\}$. Since the choice of t^\sharp was arbitrary, we have established that $\{\hat{\beta}(s) : s \in \overline{t+1, n}\} = \{\beta(l) : l \in \overline{t, n-1}\} \forall t \in \overline{1, n-1}$. But in this case, (3.84) implies that

$$\begin{aligned} & \mathbf{c}(\text{pr}_2(\hat{h}_t), \text{pr}_1(\hat{h}_{t+1}), \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) + c_{\hat{\beta}(t+1)}(\hat{h}_{t+1}, \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) = \\ & \mathbf{c}(\text{pr}_2(h_{t-1}), \text{pr}_1(h_t), \{\beta(l) : l \in \overline{t, n-1}\}) + c_{\beta(t)}(h_t, \{\beta(l) : l \in \overline{t, n-1}\}) \forall t \in \overline{1, n-1}. \end{aligned}$$

In view of that, we have the following equality:

$$\begin{aligned} & \max_{t \in \overline{1, n-1}} \left[\mathbf{c}(\text{pr}_2(\hat{h}_t), \text{pr}_1(\hat{h}_{t+1}), \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) + c_{\hat{\beta}(t+1)}(\hat{h}_{t+1}, \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) \right] = \\ & \max_{\xi \in \overline{0, n-2}} \left[\mathbf{c}(\text{pr}_2(h_\xi), \text{pr}_1(h_{\xi+1}), \{\beta(l) : l \in \overline{\xi+1, n-1}\}) + \right. \\ & \left. + c_{\beta(\xi+1)}(h_{\xi+1}, \{\beta(l) : l \in \overline{\xi+1, n-1}\}) \right]. \end{aligned} \quad (3.91)$$

From (3.43) and (3.91), we have

$$\begin{aligned} & \max_{t \in \overline{1, n-1}} \left[\mathbf{c}(\text{pr}_2(\hat{h}_t), \text{pr}_1(\hat{h}_{t+1}), \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) + c_{\hat{\beta}(t+1)}(\hat{h}_{t+1}, \{\hat{\beta}(s) : s \in \overline{t+1, n}\}) \right] = \\ & = \mathfrak{C}_Q^{(\beta)}[(h_i)_{i \in \overline{0, n-1}}]. \end{aligned}$$

Then, (3.76) implies the estimate $v(x, K) \leq \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(\mathbf{z}), K) + c_q(\mathbf{z}, K); \mathfrak{C}_Q^{(\beta)}[(h_i)_{i \in \overline{0, n-1}}] \right\} \right)$, whence, in its own turn, we obtain, in view of (3.42), the inequality

$$v(x, K) \leq \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(\mathbf{z}), K) + c_q(\mathbf{z}, K); v(\text{pr}_2(\mathbf{z}), Q) \right\} \right). \quad (3.92)$$

Now, taking into account (3.35), we obtain

$$v(x, K) \leq \min_{j \in \mathbf{I}(K)} \min_{z \in \mathbb{M}_j} \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(z), K) + c_j(z, K); v(\text{pr}_2(z), K \setminus \{j\}) \right\} \right). \quad (3.93)$$

From (3.30) and (3.93), we get the desired equality

$$v(x, K) = \min_{j \in \mathbf{I}(K)} \min_{z \in \mathbb{M}_j} \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(z), K) + c_j(z, K); v(\text{pr}_2(z), K \setminus \{j\}) \right\} \right)$$

for the case of $n \in \overline{2, N}$. In other words,

$$\begin{aligned} & (n \in \overline{2, N}) \Rightarrow \\ & \left(v(x, K) = \min_{j \in \mathbf{I}(K)} \min_{z \in \mathbb{M}_j} \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(z), K) + c_j(z, K); v(\text{pr}_2(z), K \setminus \{j\}) \right\} \right) \right). \end{aligned}$$

In view of (3.12), we have equality (3.11) in all possible cases. □

From (3.9) and Theorem 1, it follows that

$$V = \min_{j \in \mathbf{I}(\overline{1, N})} \min_{z \in \mathbb{M}_j} \sup \left(\left\{ \mathbf{c}(x^0, \text{pr}_1(z), \overline{1, N}) + c_j(z, \overline{1, N}); v(\text{pr}_2(z), \overline{1, N} \setminus \{j\}) \right\} \right). \quad (3.94)$$

4. Economical Dynamic Programming

In this section, we use the construction from [16], which stems from constructions of [13, §4.9]; however, in [13] and [16], the criterion was *additive*, whence the need to adapt the mentioned construction of [16] to the problem considered in this paper. Following [13, 16], let us introduce the family

$$\mathcal{G} \triangleq \left\{ K \in \mathfrak{N} \mid \forall z \in \mathbf{K} (\text{pr}_1(z) \in K) \Rightarrow (\text{pr}_2(z) \in K) \right\}, \quad (4.1)$$

elements of which (finite subsets of $\overline{1, N}$) are called *feasible (task) lists*. Assume $\mathcal{G}_s \triangleq \left\{ K \in \mathcal{G} \mid s = |K| \right\} \forall s \in \overline{1, N}$. The tuple $(\mathcal{G}_s)_{s \in \overline{1, N}}$ evidently defines an (ordered) partition of \mathcal{G} (4.1). Under this partition, $\mathcal{G}_N = \{\overline{1, N}\}$ (the singleton reflecting the *complete* task list) and $\mathcal{G}_1 = \{\{t\} : t \in \overline{1, N} \setminus \{\mathbf{K}_1\}\}$, where $\mathbf{K}_1 = \{\text{pr}_1(z) : z \in \mathbf{K}\}$; thus, \mathcal{G}_1 contains singletons corresponding to "nonsenders". In addition,

$$\mathcal{G}_{s-1} = \left\{ K \setminus \{t\} : K \in \mathcal{G}_s, t \in \mathbf{I}(K) \right\} \quad \forall s \in \overline{2, N}; \quad (4.2)$$

see [16].

In view of (4.2), we obtain the recurrence relation $\mathcal{G}_N \longrightarrow \mathcal{G}_{N-1} \longrightarrow \dots \longrightarrow \mathcal{G}_1$, which provides for construction of family (4.1) (recall that \mathcal{G}_N is known). It is also known [16] that $\mathcal{G}_s \neq \emptyset \quad \forall s \in \overline{1, N}$. If $s \in \overline{1, N-1}$ and $K \in \mathcal{G}_s$, then [16], in terms of the set

$$\mathcal{J}_s(K) \triangleq \{j \in \overline{1, N} \setminus K \mid \{j\} \cup K \in \mathcal{G}_{s+1}\} \in \mathcal{P}'(\overline{1, N} \setminus K), \quad (4.3)$$

construct $\mathcal{M}_s[K]$ in the form of union of all \mathbf{M}_j , $j \in \mathcal{J}_s(K)$; let us also construct a *cell* $\mathbb{D}_s[K] \triangleq \{(x, K) : x \in \mathcal{M}_s[K]\}$ of the *position* space. Then, as in [16], assume

$$D_s \triangleq \bigcup_{K \in \mathcal{G}_s} \mathbb{D}_s[K] \in \mathcal{P}'(\mathbf{X} \times \mathcal{G}_s) \quad \forall s \in \overline{1, N-1}; \quad (4.4)$$

we have defined (intermediate) *layers* of the position space. Let us also define the two (nonempty) "border" layers: $D_0 \triangleq \{(x, \emptyset) : x \in \mathfrak{M}\}$, where \mathfrak{M} is the union of all the sets \mathbf{M}_j , $j \in \overline{1, N} \setminus \mathbf{K}_1$ and $D_N \triangleq \{(x^0, \overline{1, N})\}$ (the singleton that contains the position $(x^0, \overline{1, N})$). The tuple $(D_s)_{s \in \overline{0, N}}$ possesses the following important property (see [13, §4.9]):

$$(y, K \setminus \{k\}) \in D_{s-1} \quad \forall s \in \overline{1, N} \quad \forall K \in \mathcal{G}_s \quad \forall k \in \mathbf{I}(K) \quad \forall y \in \mathbf{M}_k. \quad (4.5)$$

In connection with (4.5), recall that, according to (2.5), $\text{pr}_2(z) \in \mathbf{M}_k \quad \forall k \in \overline{1, N} \quad \forall z \in \mathbb{M}_k$. In addition, (4.4) implies that, for $s \in \overline{1, N-1}$ and $(x, K) \in D_s$, we definitely have $K \in \mathcal{G}_s$ (by definition of *cells* of the space of *positions*); in case $(x, K) \in D_N$, we have $K = \overline{1, N} \in \mathcal{G}_N$. Thus, $K \in \mathcal{G}_s \quad \forall s \in \overline{1, N} \quad \forall (x, K) \in D_s$. Taking into account (4.5), we consequently obtain the property

$$(\text{pr}_2(z), K \setminus \{k\}) \in D_{s-1} \quad \forall s \in \overline{1, N} \quad \forall (x, K) \in D_s \quad \forall k \in \mathbf{I}(K) \quad \forall z \in \mathbb{M}_k. \quad (4.6)$$

Now, bearing in mind the fact that $D_0 \in \mathcal{P}'(\mathbf{X} \times \mathcal{P}(\overline{1, N}))$, $D_1 \in \mathcal{P}'(\mathbf{X} \times \mathcal{P}(\overline{1, N}))$, \dots , and $D_N \in \mathcal{P}'(\mathbf{X} \times \mathcal{P}(\overline{1, N}))$ (we also have to consider $\mathfrak{N} \subset \mathcal{P}(\overline{1, N})$), we define the *restrictions* of the Bellman function onto the layers of the space of positions. Namely, for $s \in \overline{0, N}$, we assume $v_s \in \mathcal{R}_+[D_s]$ to be a function such that $v_s(x, K) \triangleq v(x, K) \quad \forall (x, K) \in D_s$. Property (4.6) now implies that, for $s \in \overline{1, N}$ and $(x, K) \in D_s$, we have $v_{s-1}(\text{pr}_2(h), K \setminus \{k\}) \in [0, \infty[\quad \forall k \in \mathbf{I}(K) \quad \forall h \in \mathbb{M}_k$; hence, we can determine the value

$$\min_{j \in \mathbf{I}(K)} \min_{z \in \mathbb{M}_j} \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(z), K) + c_j(z, K); v_{s-1}(\text{pr}_2(z), K \setminus \{j\}) \right\} \right) \in [0, \infty[.$$

Now, Theorem 1 implies, in view of (4.5), the following proposition:

Proposition 1. *If $s \in \overline{1, N}$, then the function v_s is obtained from the function v_{s-1} by the rule*

$$v_s(x, K) = \min_{j \in \mathbf{I}(K)} \min_{z \in \mathbb{M}_j} \sup \left(\left\{ \mathbf{c}(x, \text{pr}_1(z), K) + c_j(z, K); v_{s-1}(\text{pr}_2(z), K \setminus \{j\}) \right\} \right) \quad \forall (x, K) \in D_s.$$

Note that the function v_0 is known: $v_0(x, \emptyset) = 0 \quad \forall x \in \mathfrak{M}$. Therefore, Proposition 1 defines the recurrent procedure

$$v_0 \longrightarrow v_1 \longrightarrow \dots \longrightarrow v_N. \quad (4.7)$$

Implementation of (4.7) would yield (see (3.94)) the value of problem (2.11):

$$V = v_N(x^0, \overline{1, N}) = \min_{j \in \mathbf{I}(\overline{1, N})} \min_{z \in \mathbb{M}_j} \sup \left(\left\{ \mathbf{c}(x^0, \text{pr}_1(z), \overline{1, N}) + c_j(z, \overline{1, N}); v_{N-1}(\text{pr}_2(z), \overline{1, N} \setminus \{j\}) \right\} \right). \quad (4.8)$$

This value may be used to check if movements (2.2) are possible. Namely, we might be given a *tolerance threshold* \mathbf{d} of harmful effects, effective on each step of process (2.2). Indeed, if $V \leq \mathbf{d}$, then, in compliance with (2.11) and (2.12), we may organize the motion (2.2) such that

$$\begin{aligned} & \mathbf{c}(\text{pr}_2(z^{(t)}), \text{pr}_1(z^{(t+1)}), \{\alpha(s) : s \in \overline{t+1, N}\}) + \\ & + c_{\alpha(t+1)}(z^{(t+1)}, \{\alpha(s) : s \in \overline{t+1, N}\}) \leq \mathbf{d} \quad \forall s \in \overline{0, N-1}. \end{aligned} \quad (4.9)$$

However, in case $\mathbf{d} < V$, it is not possible to provide (4.9). Thus, the implementation of (4.7), which is completely defined by Proposition 1, already yields useful data.

Optimal solutions are also constructed based on Proposition 1. Assume $\mathbf{z}^{(0)} \triangleq (x^0, x^0)$; this is an OP from $\mathbb{X} \times \mathbf{X}$. Then, using (4.8), choose $j_1 \in \mathbf{I}(\overline{1, N})$ and $\mathbf{z}^{(1)} \in \mathbb{M}_{j_1}$ such that

$$V = \sup \left(\left\{ \mathbf{c}(x^0, \text{pr}_1(\mathbf{z}^{(1)}), \overline{1, N}) + c_{j_1}(\mathbf{z}^{(1)}, \overline{1, N}); v_{N-1}(\text{pr}_2(\mathbf{z}^{(1)}), \overline{1, N} \setminus \{j_1\}) \right\} \right) \quad (4.10)$$

(find the minimum in (4.8)). By (4.6), we have (since $(x^0, \overline{1, N}) \in D_N$)

$$(\text{pr}_2(\mathbf{z}^{(1)}), \overline{1, N} \setminus \{j_1\}) \in D_{N-1}; \quad (4.11)$$

hence, Proposition 1 implies that

$$\begin{aligned} & v_{N-1}(\text{pr}_2(\mathbf{z}^{(1)}), \overline{1, N} \setminus \{j_1\}) = \\ & \min_{j \in \mathbf{I}(\overline{1, N} \setminus \{j_1\})} \min_{z \in \mathbb{M}_j} \sup \left(\left\{ \mathbf{c}(\text{pr}_2(\mathbf{z}^{(1)}), \text{pr}_1(z), \overline{1, N} \setminus \{j_1\}) + \right. \right. \\ & \left. \left. + c_j(z, \overline{1, N} \setminus \{j_1\}); v_{N-2}(\text{pr}_2(z), \overline{1, N} \setminus \{j_1, j\}) \right\} \right) \end{aligned} \quad (4.12)$$

(take into account $\overline{1, N} \setminus \{j_1, j\} = (\overline{1, N} \setminus \{j_1\}) \setminus \{j\}$ for $j \in \mathbf{I}(\overline{1, N} \setminus \{j_1\})$). Now, find the minimum in (4.12): find $j_2 \in \mathbf{I}(\overline{1, N} \setminus \{j_1\})$ and $\mathbf{z}^{(2)} \in \mathbb{M}_{j_2}$ such that

$$\begin{aligned} & v_{N-1}(\text{pr}_2(\mathbf{z}^{(1)}), \overline{1, N} \setminus \{j_1\}) = \\ & \sup \left(\left\{ \mathbf{c}(\text{pr}_2(\mathbf{z}^{(1)}), \text{pr}_1(\mathbf{z}^{(2)}), \overline{1, N} \setminus \{j_1\}) + c_{j_2}(\mathbf{z}^{(2)}, \overline{1, N} \setminus \{j_1\}); \right. \right. \\ & \left. \left. v_{N-2}(\text{pr}_2(\mathbf{z}^{(2)}), \overline{1, N} \setminus \{j_1, j_2\}) \right\} \right), \end{aligned} \quad (4.13)$$

where (see (4.6)) the property

$$(\text{pr}_2(\mathbf{z}^{(2)}), \overline{1, N} \setminus \{j_1, j_2\}) = (\text{pr}_2(\mathbf{z}^{(2)}), \overline{1, N} \setminus \{j_k : k \in \overline{1, 2}\}) \in D_{N-2} \quad (4.14)$$

holds. Expressions (4.10) and (4.13) imply that, in particular,

the parallelization was made with the aid of shared memory multiprocessing API OpenMP 3.0. The calculation took 1h. 46min. 34sec. on 12 cores.

Conclusion

The paper proposes an unorthodox variant of the dynamic programming method to solve precedence-constrained routing problems with nonadditive cost aggregation (the bottleneck cost aggregation) and dependence on the list of pending tasks. A method of obtaining the optimal solutions is specified, and the (optimal) algorithm is implemented in the form of a parallel computer program.

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МОДЕЛЬ «НЕАДДИТИВНОЙ» ЗАДАЧИ МАРШРУТИЗАЦИИ С ФУНКЦИЯМИ СТОИМОСТИ, ЗАВИСЯЩИМИ ОТ СПИСКА ЗАДАНИЙ

А.Г. Ченцов, Я.В. Салий

Рассматривается следующий (усложненный) вариант маршрутной задачи «на узкие места»: исследуется задача последовательного обхода мегаполисов, осложненная условиями предшествования и тем, что функции стоимости (перемещений и внутренних работ) могут явным образом зависеть от списка заданий, которые не выполнены на данный момент. Процесс перемещений рассматривается в виде совокупности этапов, включающих внешнее перемещение к соответствующему мегаполису и последующее выполнение (внутренних по смыслу) работ, связанных с данным мегаполисом. Качество совокупного процесса оценивается максимумом стоимостей составляющих его этапов; рассматривается задача на минимум упомянутого критерия (получается задача на минимум, обычно именуемая задачей «на узкие места»). Для построения оптимального решения в виде пары маршрут-трасса (трасса, или траектория, соответствует конкретному варианту прохождения мегаполисов, нумеруемых в соответствии с маршрутом, определяемым в виде перестановки индексов) построен «нестандартный» вариант метода динамического программирования, при реализации которого не используется, в случае ограничений в виде условий предшествования, построение всего массива значений функции Беллмана.

Ключевые слова: динамическое программирование; маршрут; условия предшествования.

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